



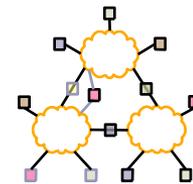
# CE693: Adv. Computer Networking

## L-09 Wireless in the Real World

*Acknowledgments: Lecture slides are from the graduate level Computer Networks course taught by Srinivasan Seshan at CMU. When slides are obtained from other sources, a reference will be noted on the bottom of that slide. A full list of references is provided on the last slide.*

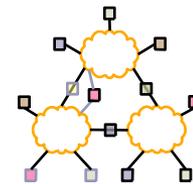


# Wireless in the Real World



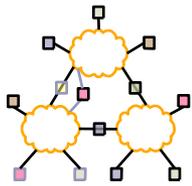
- Real world deployment patterns
- Mesh networks and deployments
- Assigned reading
  - Architecture and Evaluation of an Unplanned 802.11b Mesh Network
  - White Space Networking with Wi-Fi like Connectivity

# Wireless Challenges



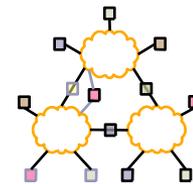
- Force us to rethink many assumptions
- Need to share airwaves rather than wire
  - Don't know what hosts are involved
  - Host may not be using same link technology
- Mobility
- Other characteristics of wireless
  - Noisy → lots of losses
  - Slow
  - Interaction of multiple transmitters at receiver
    - Collisions, capture, interference
  - Multipath interference

# Overview



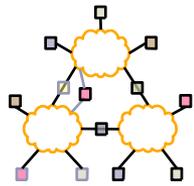
- 802.11
  - Deployment patterns
  - Reaction to interference
  - Interference mitigation
- Mesh networks
  - Architecture
  - Measurements
- White space networks

# Characterizing Current Deployments



- Datasets
- Place Lab: 28,000 APs
  - MAC, ESSID, GPS
  - Selected US cities
  - [www.placelab.org](http://www.placelab.org)
- Wifimaps: 300,000 APs
  - MAC, ESSID, Channel, GPS (derived)
  - [wifimaps.com](http://wifimaps.com)
- Pittsburgh Wardrive: 667 APs
  - MAC, ESSID, Channel, Supported Rates, GPS

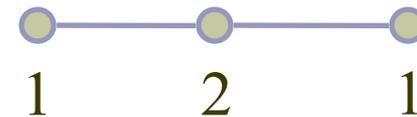
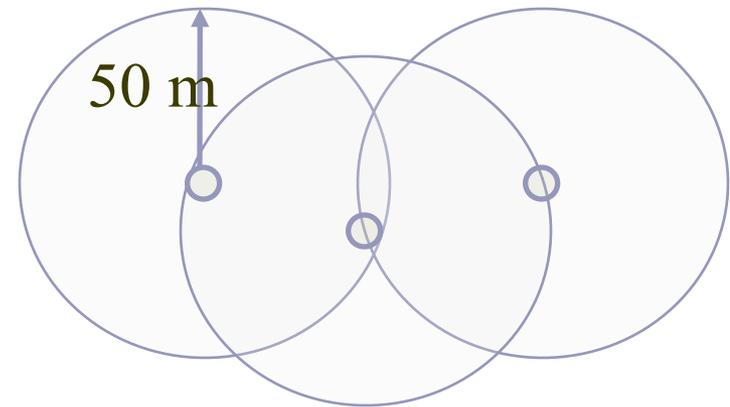
# AP Stats, Degrees: Placelab



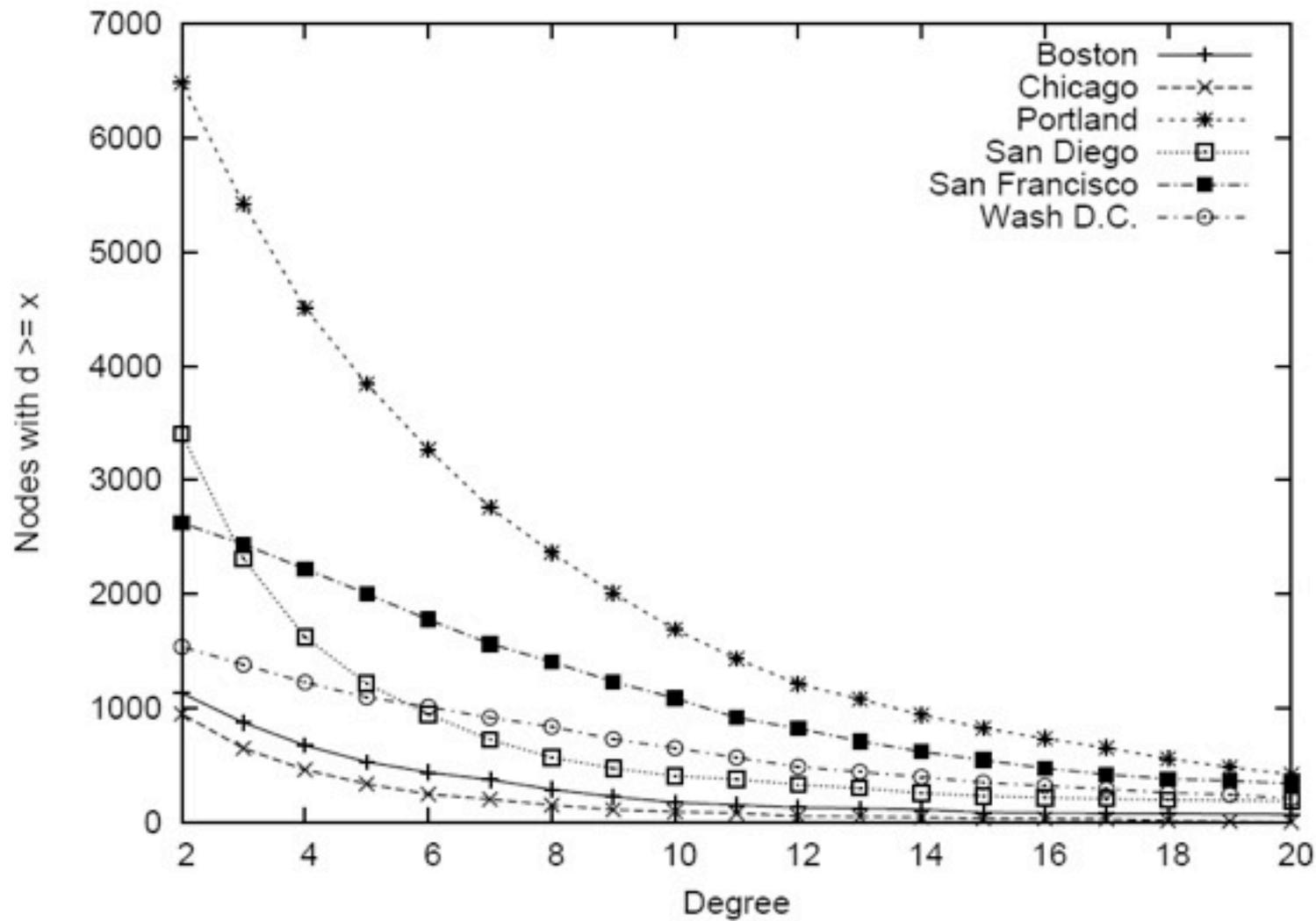
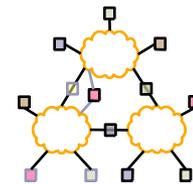
(Placelab: 28000 APs, MAC, ESSID, GPS)

#APs Max.  
degree

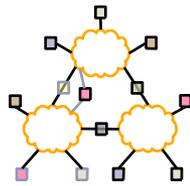
<b>Portland</b>	8683	54
<b>San Diego</b>	7934	76
<b>San Francisco</b>	3037	85
<b>Boston</b>	2551	39



# Degree Distribution: Place Lab



# Unmanaged Devices



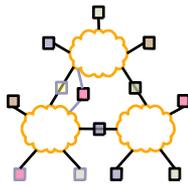
WifiMaps.com  
(300,000 APs, MAC, ESSID, Channel)

## Channel %age

<b>6</b>	51
<b>11</b>	21
<b>1</b>	14
<b>10</b>	4

- Most users don't change default channel
- Channel selection must be automated

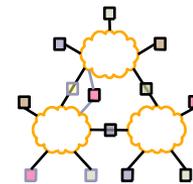
# Growing Interference in Unlicensed Bands



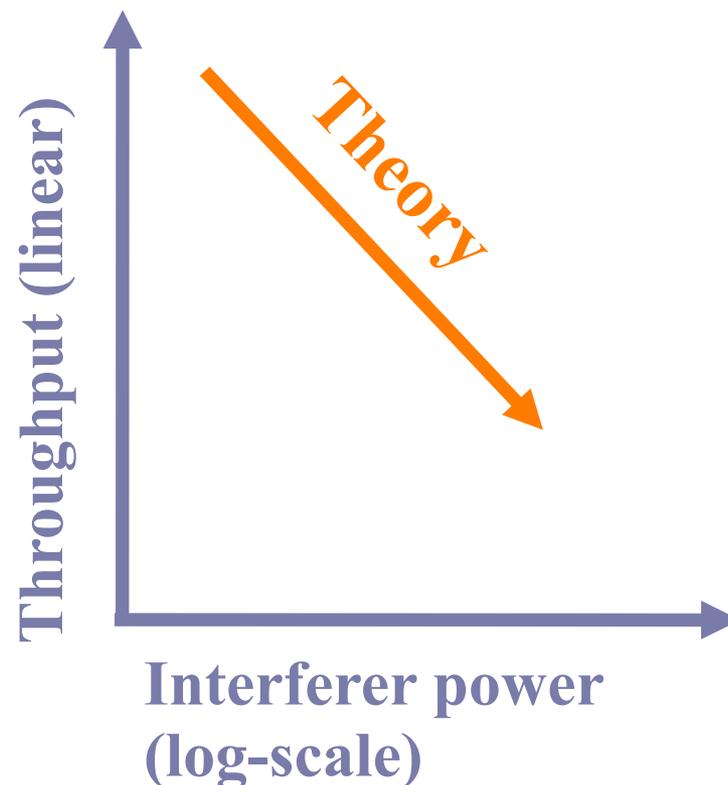
- Anecdotal evidence of problems, but how severe?
- Characterize how 802.11 operates under interference in practice



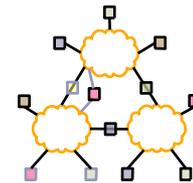
# What do we expect?



- Throughput to decrease linearly with interference
- There to be lots of options for 802.11 devices to tolerate interference
  - Bit-rate adaptation
  - Power control
  - FEC
  - Packet size variation
  - Spread-spectrum processing
  - Transmission and reception diversity

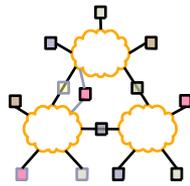


# Key Questions

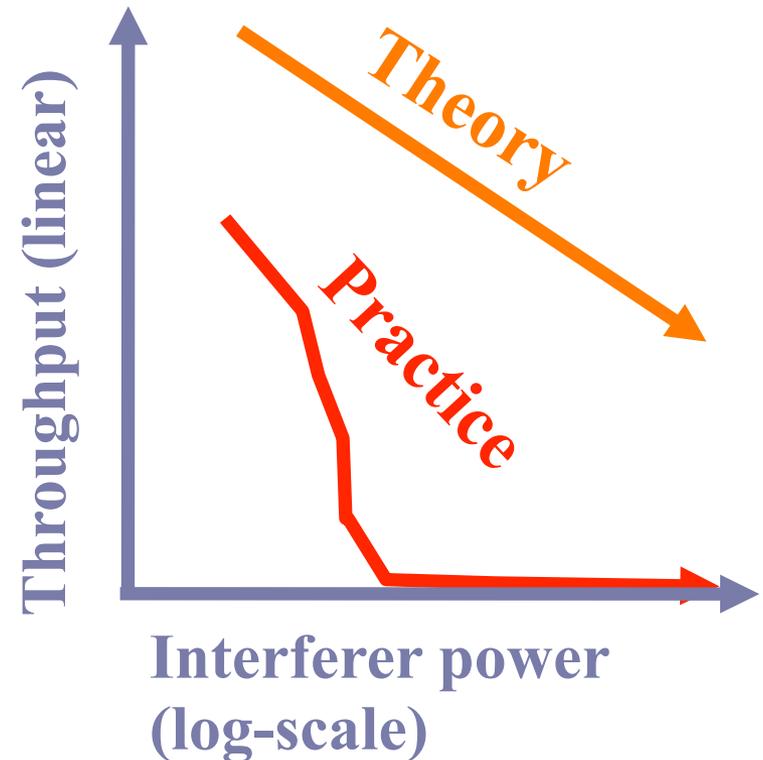


- How damaging can a low-power and/or narrow-band interferer be?
- How can today's hardware tolerate interference well?
  - What 802.11 options work well, and why?

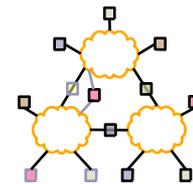
# What we see



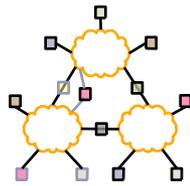
- Effects of interference more severe in practice
- Caused by hardware limitations of commodity cards, which theory doesn't model



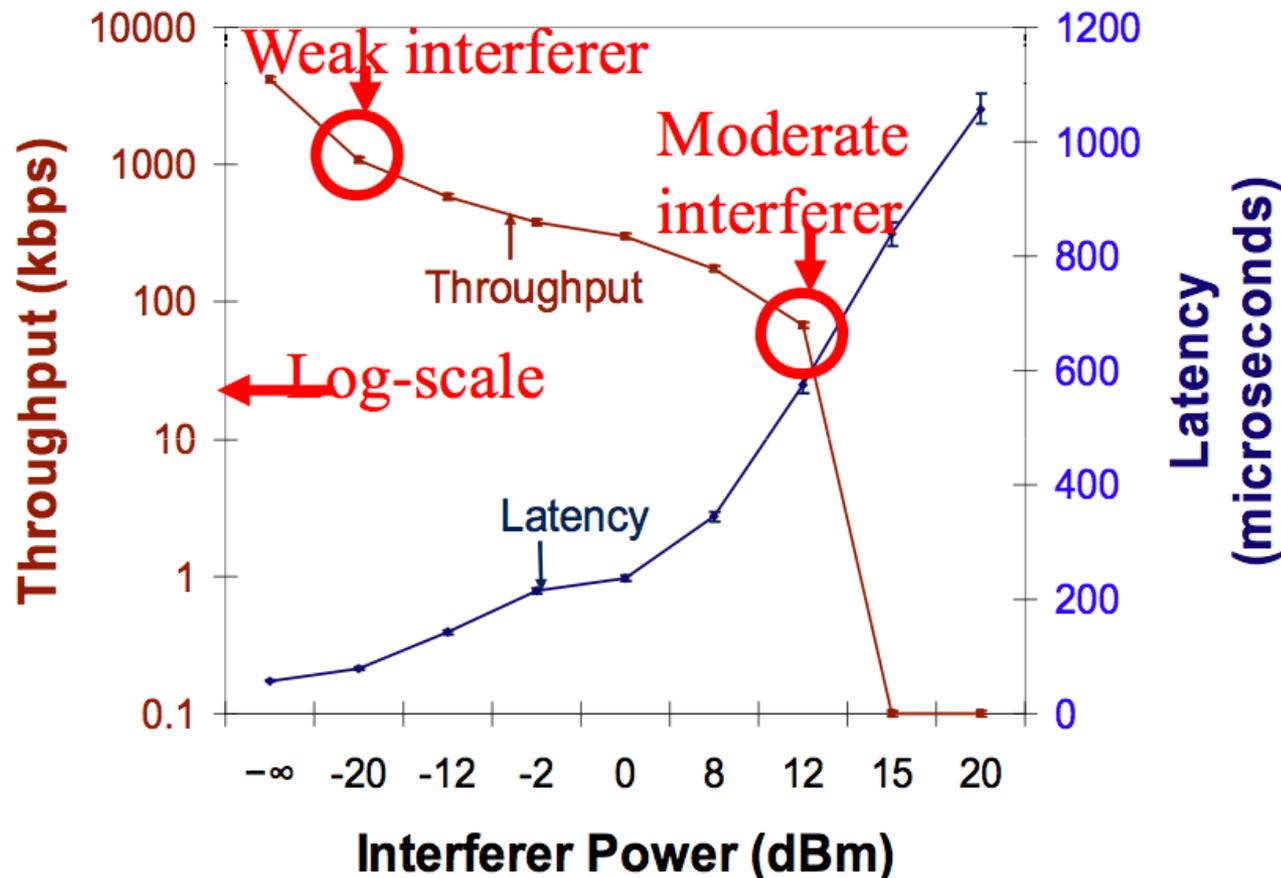
# Experimental Setup



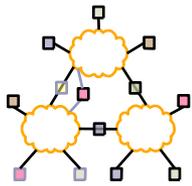
# Timing Recovery Interference



- Interferer sends continuous SYNC pattern
- Interferes with packet acquisition (PHY reception errors)

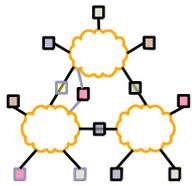


# Interference Management



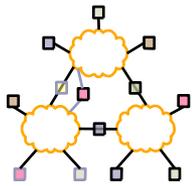
- Interference will get worse
  - Density/device diversity is increasing
  - Unlicensed spectrum is not keeping up
- Spectrum management
  - “Channel hopping” 802.11 effective at mitigating some performance problems [Sigcomm07]
  - Coordinated spectrum use – based on RF sensor network
- Transmission power control
  - Enable spatial reuse of spectrum by controlling transmit power
  - Must also adapt carrier sense behavior to take advantage

# Overview



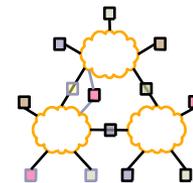
- 802.11
  - Deployment patterns
  - Reaction to interference
  - Interference mitigation
- **Mesh networks**
  - **Architecture**
  - **Measurements**
- White space networks

# Roofnet



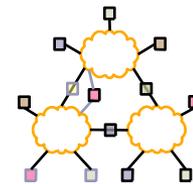
- Share a few wired Internet connections
- Goals
  - Operate without extensive planning or central management
  - Provide wide coverage and acceptable performance
- Design decisions
  - Unconstrained node placement
  - Omni-directional antennas
  - Multi-hop routing
  - Optimization of routing for throughput in a slowly changing network

# Roofnet Design



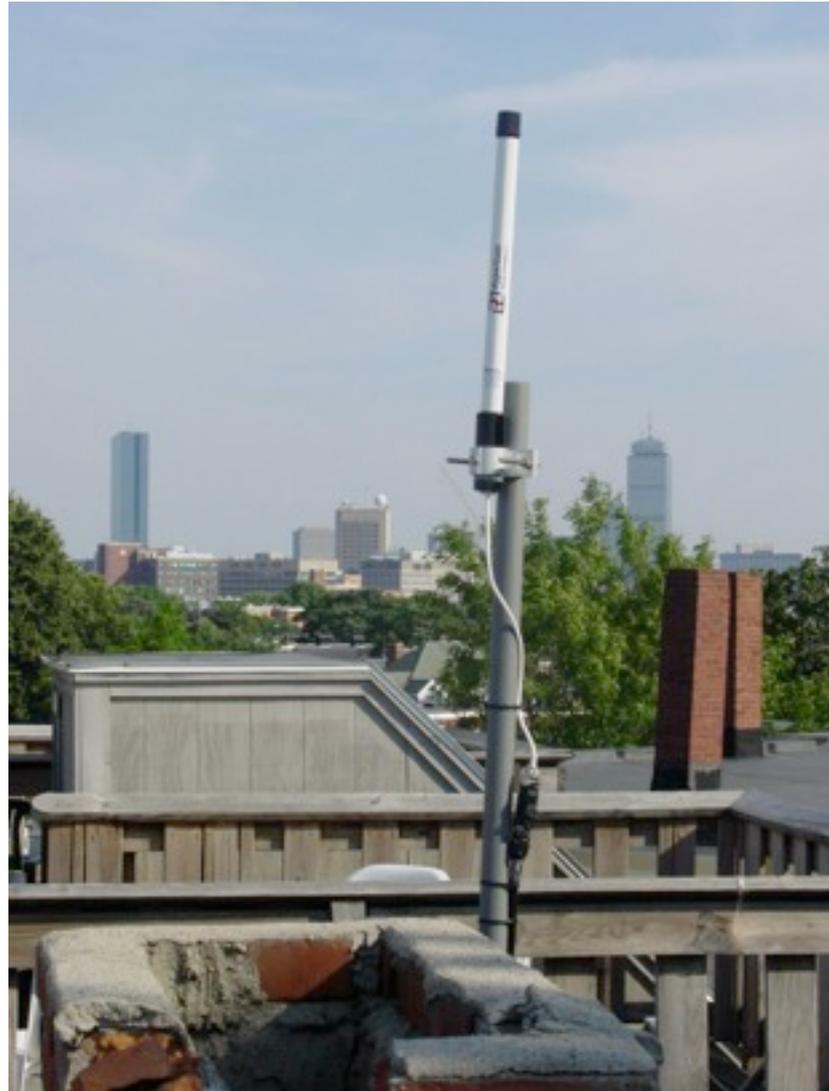
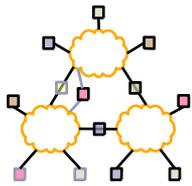
- Deployment
  - Over an area of about four square kilometers in Cambridge, Massachusetts
  - Most nodes are located in buildings
    - 3~4 story apartment buildings
    - 8 nodes are in taller buildings
  - Each Roofnet node is hosted by a volunteer user
- Hardware
  - PC, omni-directional antenna, hard drive ...
  - 802.11b card
    - RTS/CTS disabled
    - Share the same 802.11b channel
    - Non-standard “pseudo-IBSS” mode
      - Similar to standard 802.11b IBSS (ad hoc)
      - Omit beacon and BSSID (network ID)

# Roofnet Node Map

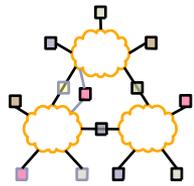


1 kilometer  
↔

# Typical Rooftop View



# A Roofnet Self-Installation Kit



## Antenna (\$65)

8dBi, 20 degree vertical

## Computer (\$340)

533 MHz PC, hard disk, CDROM

## 802.11b card (\$155)

Engenius Prism 2.5, 200mW



## 50 ft. Cable (\$40)

Low loss (3dB/100ft)

## Miscellaneous (\$75)

Chimney Mount, Lightning Arrestor, etc.

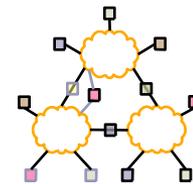
## Software (“free”)

Our networking software based on Click

**Total: \$685**

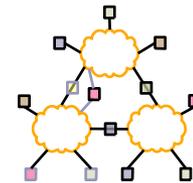
**Takes a user about 45 minutes to install on a flat roof**

# Software and Auto-Configuration



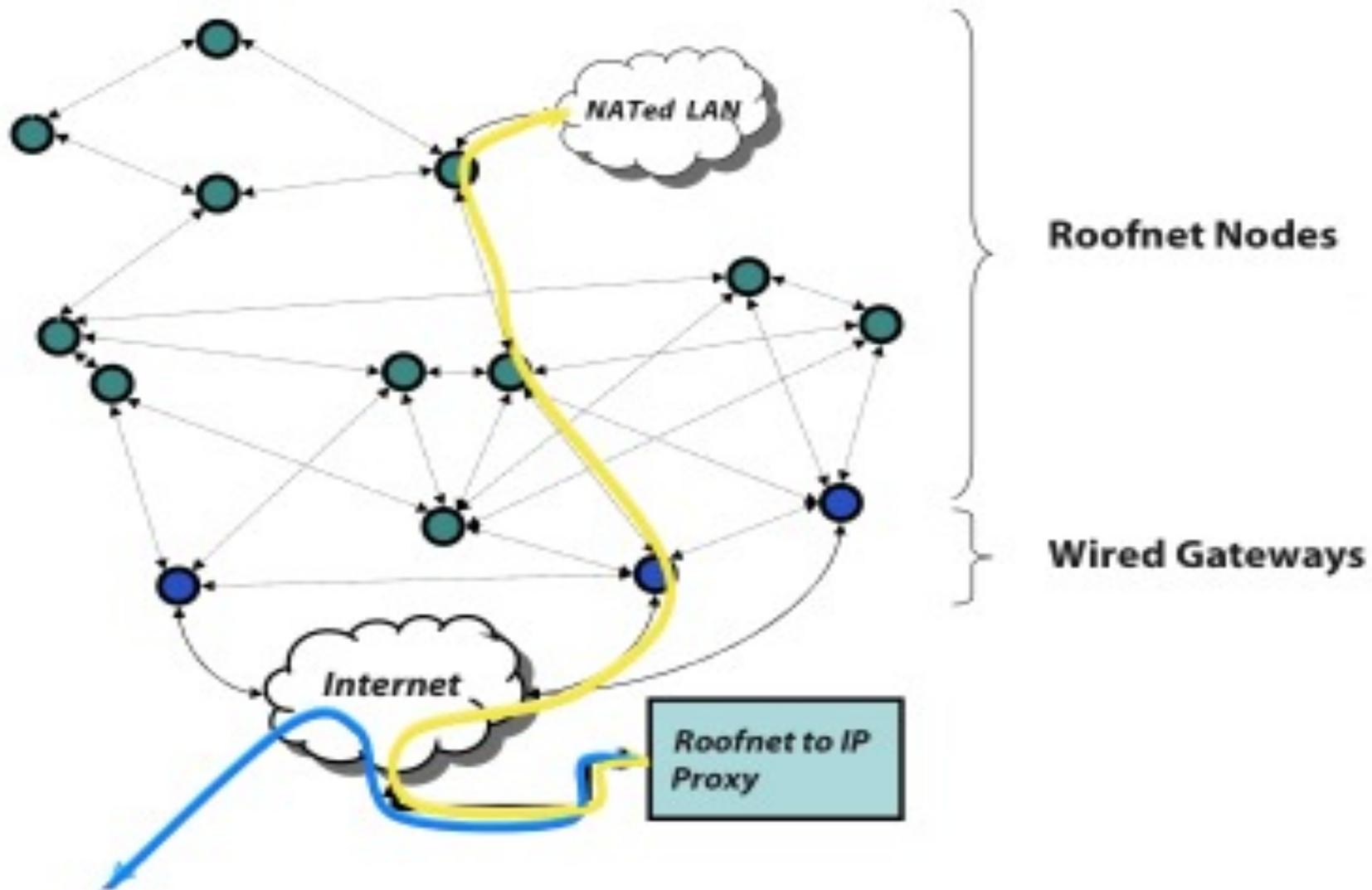
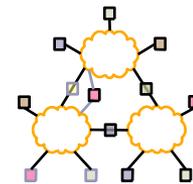
- Linux, routing software, DHCP server, web server ...
- Automatically solve a number of problems
  - Allocating addresses
  - Finding a gateway between Roofnet and the Internet
  - Choosing a good multi-hop route to that gateway
- Addressing
  - Roofnet carries IP packets inside its own header format and routing protocol
  - Assign addresses automatically
  - Only meaningful inside Roofnet, not globally routable
  - The address of Roofnet nodes
    - Low 24 bits are the low 24 bits of the node's Ethernet address
    - High 8 bits are an unused class-A IP address block
  - The address of hosts
    - Allocate 192.168.1.x via DHCP and use NAT between the Ethernet and Roofnet

# Software and Auto-Configuration

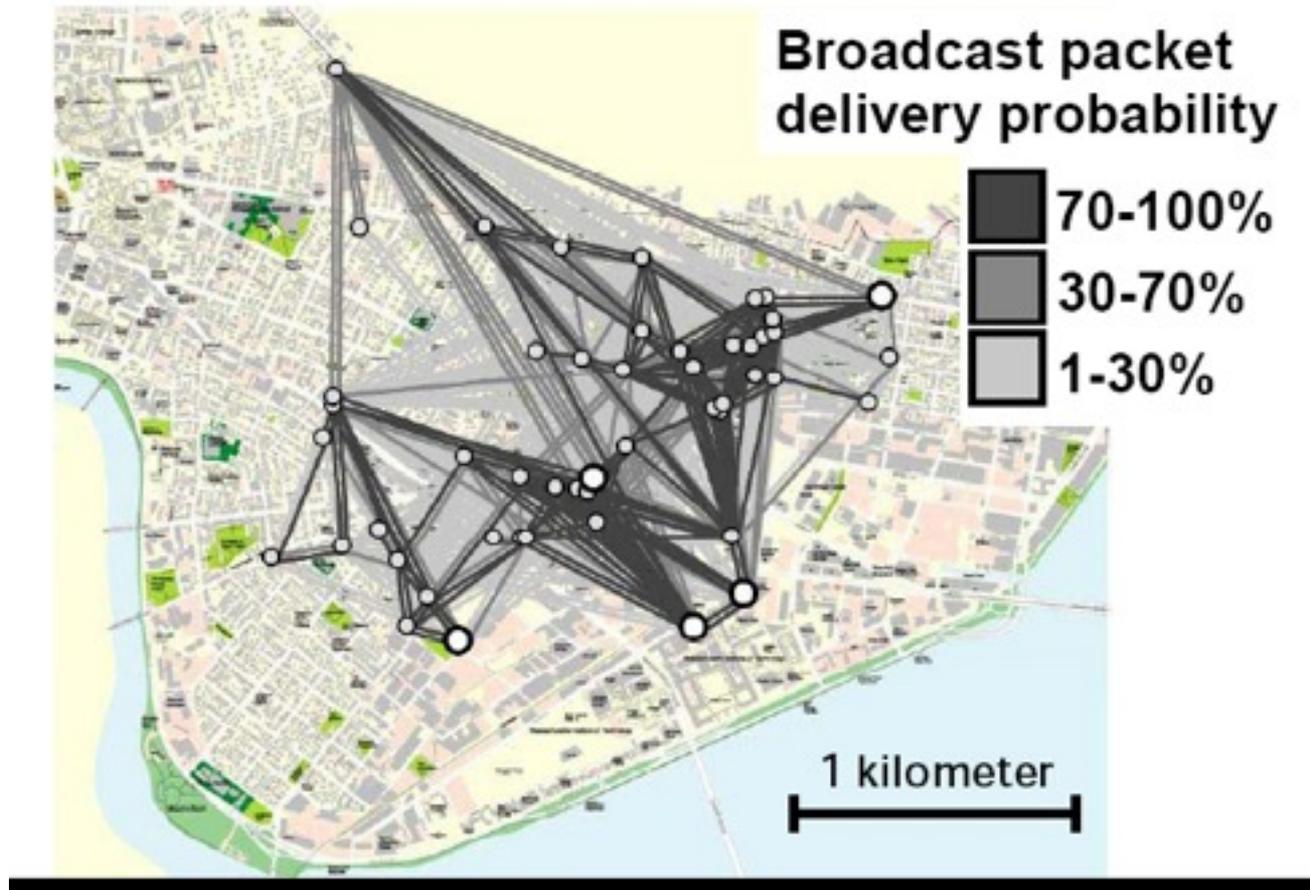
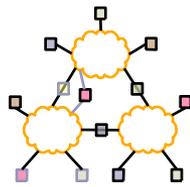


- Gateway and Internet Access
  - A small fraction of Roofnet users will share their wired Internet access links
  - Nodes which can reach the Internet
    - Advertise itself to Roofnet as an Internet gateway
    - Acts as a NAT for connection from Roofnet to the Internet
  - Other nodes
    - Select the gateway which has the best route metric
  - Roofnet currently has four Internet gateways

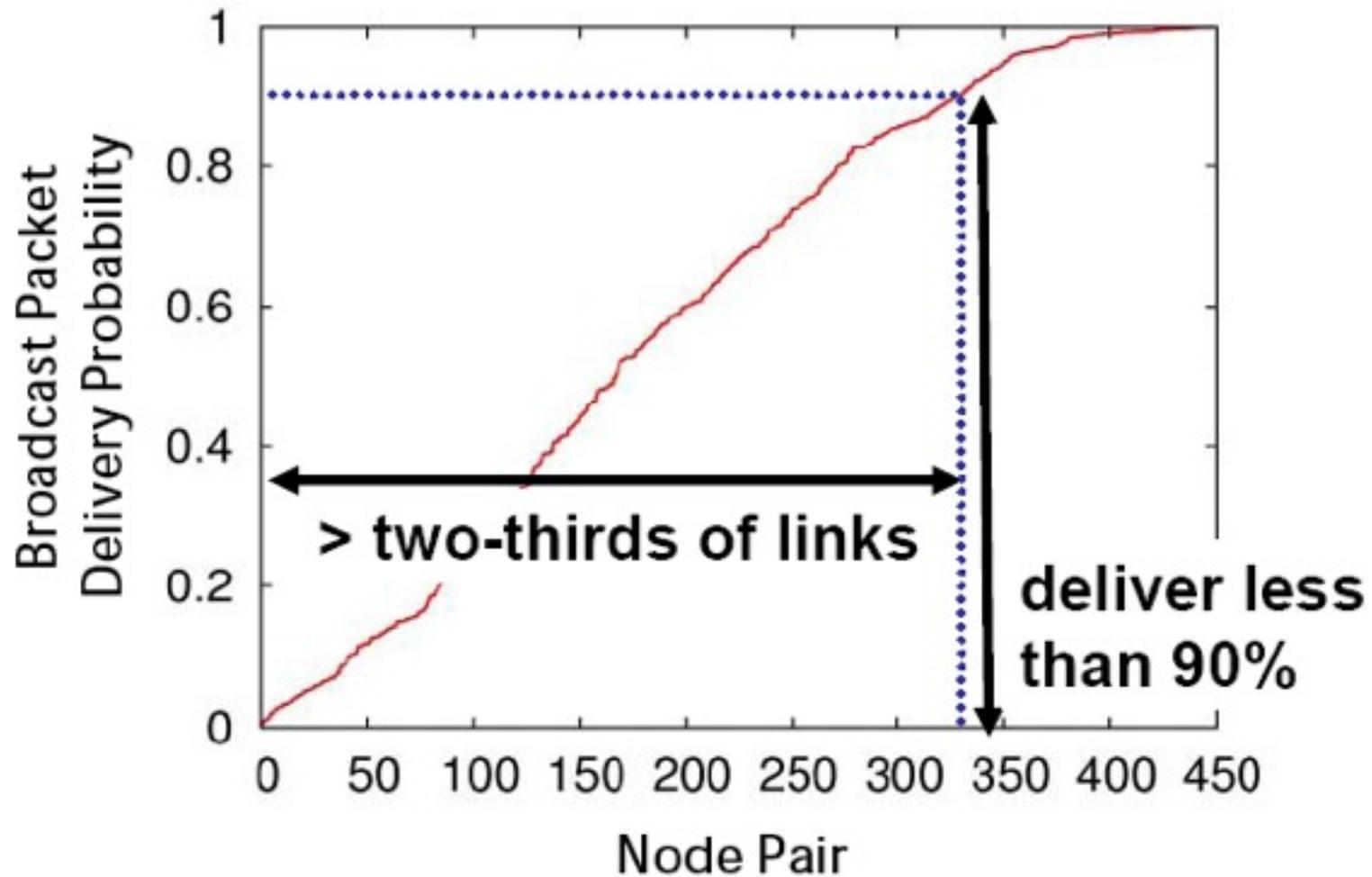
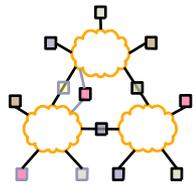
# Roofnet

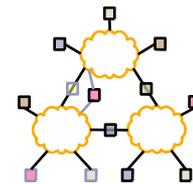


# Lossy Links are Common

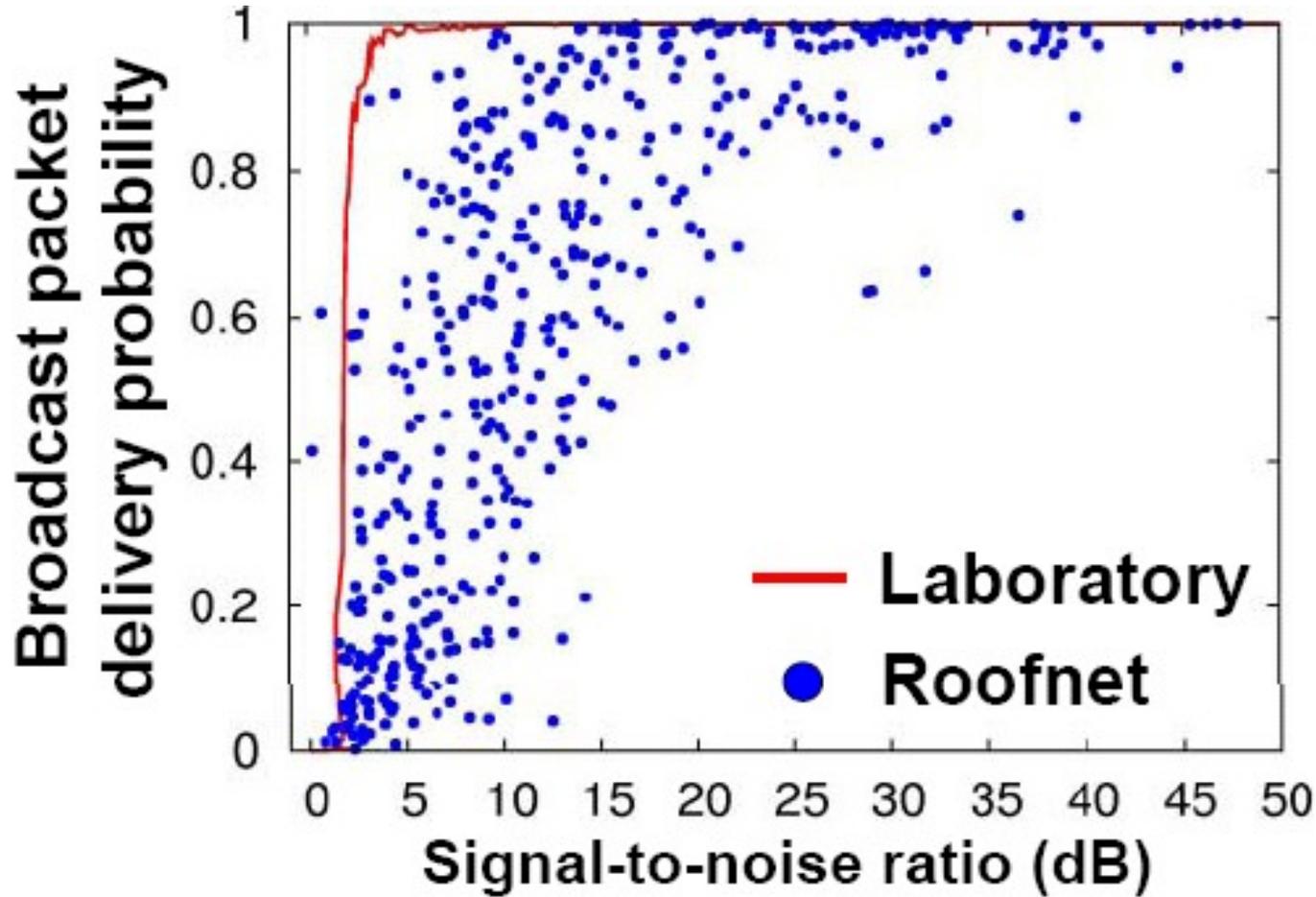


# Delivery Probabilities are Uniformly Distributed

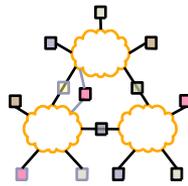




# Delivery vs. SNR

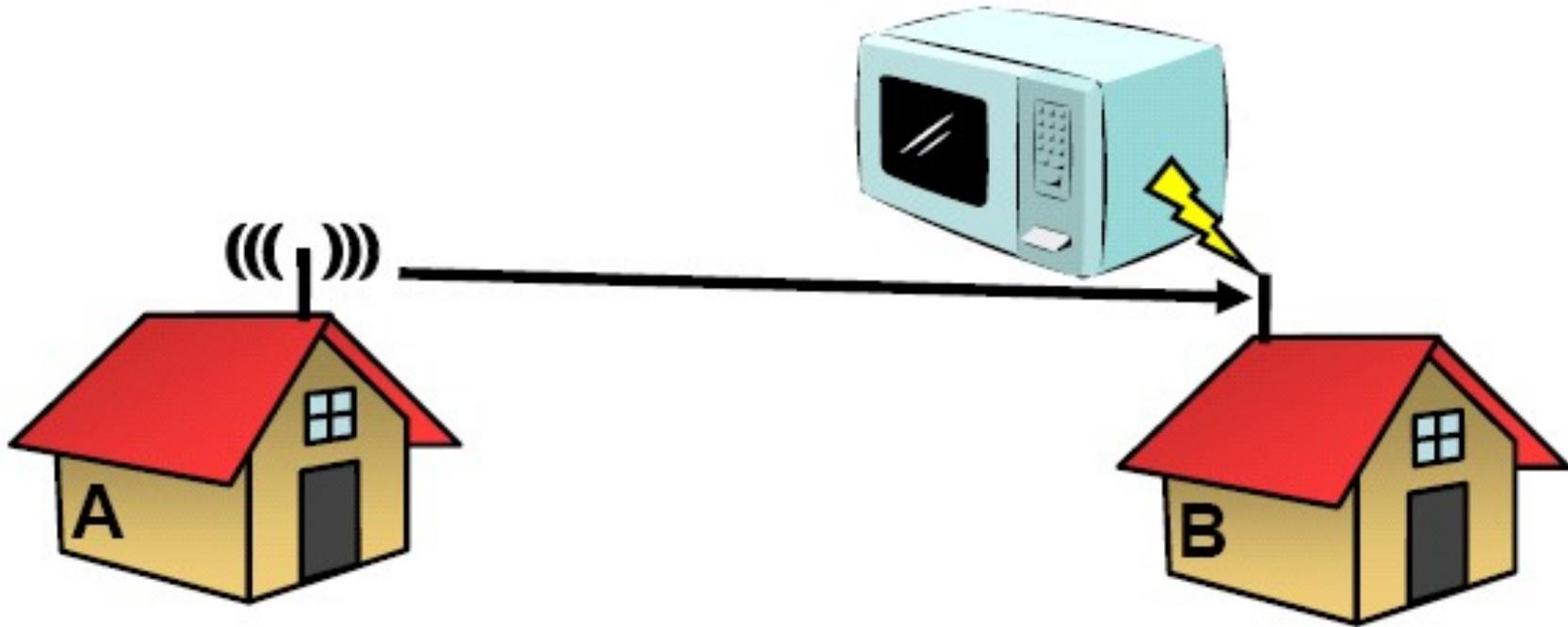


- SNR not a good predictor

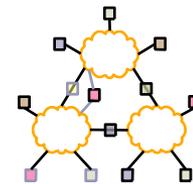


# Is it Bursty Interference?

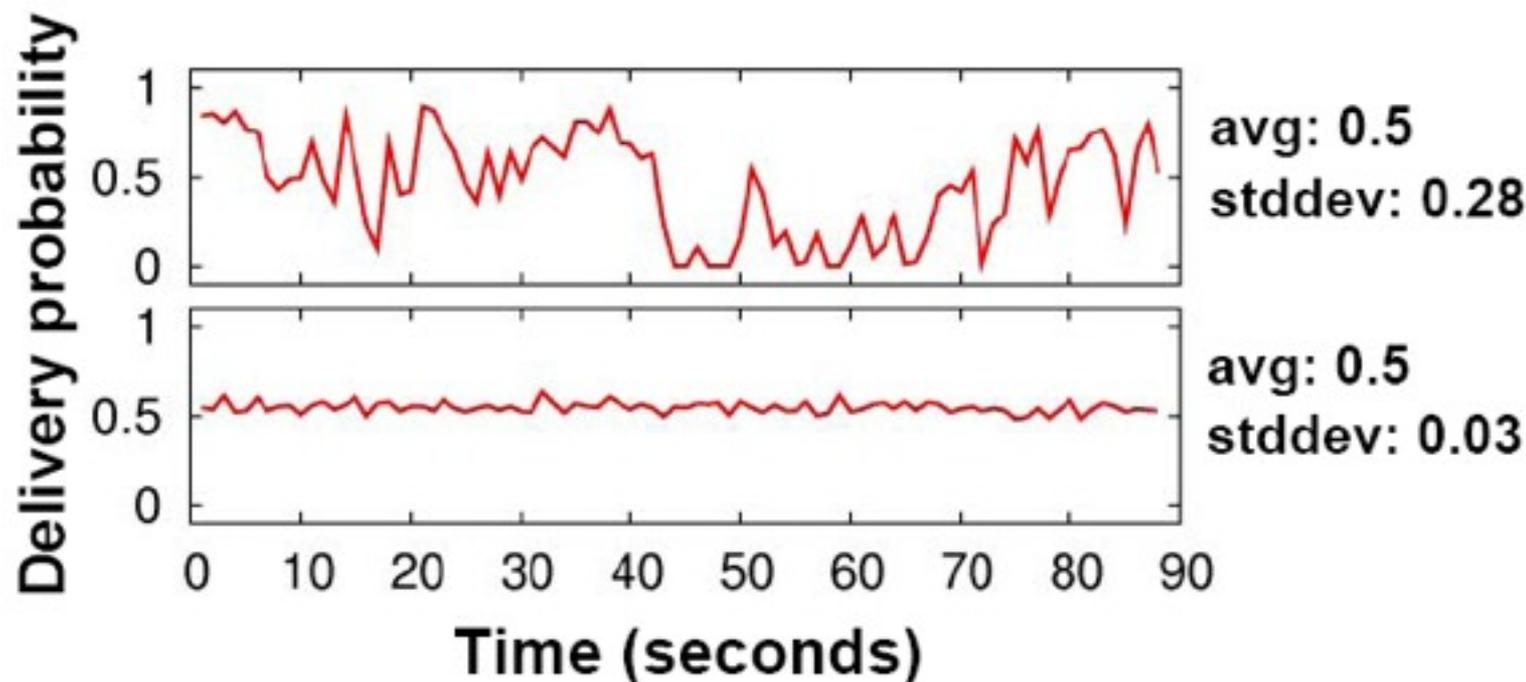
- May interfere but not impact SNR measurement



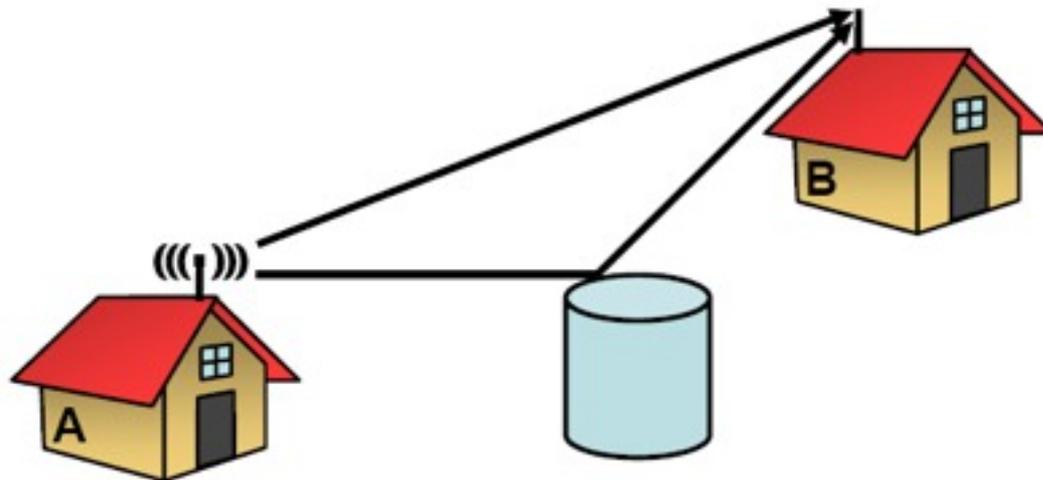
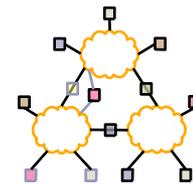
# Two Different Roofnet Links



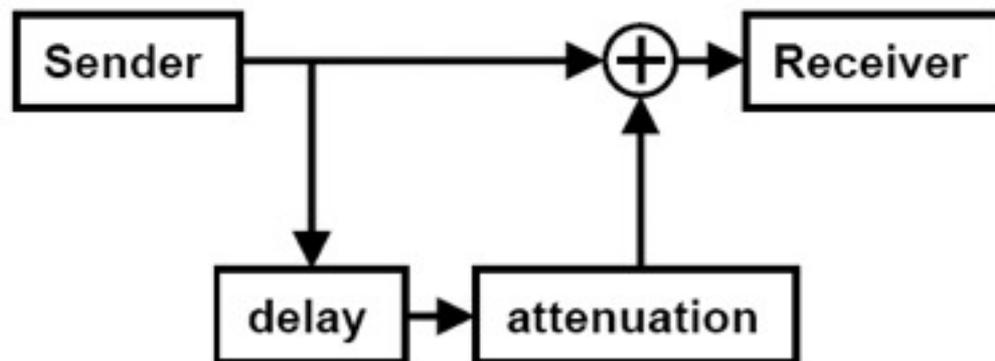
- Top is typical of bursty interference, bottom is not
- Most links are like the bottom



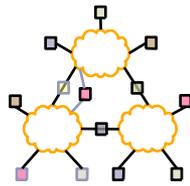
# Is it Multipath Interference?



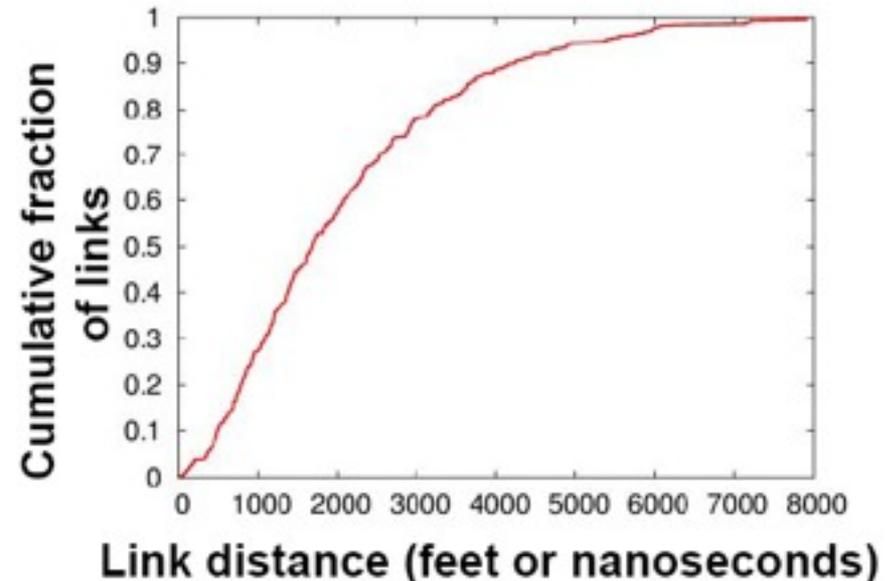
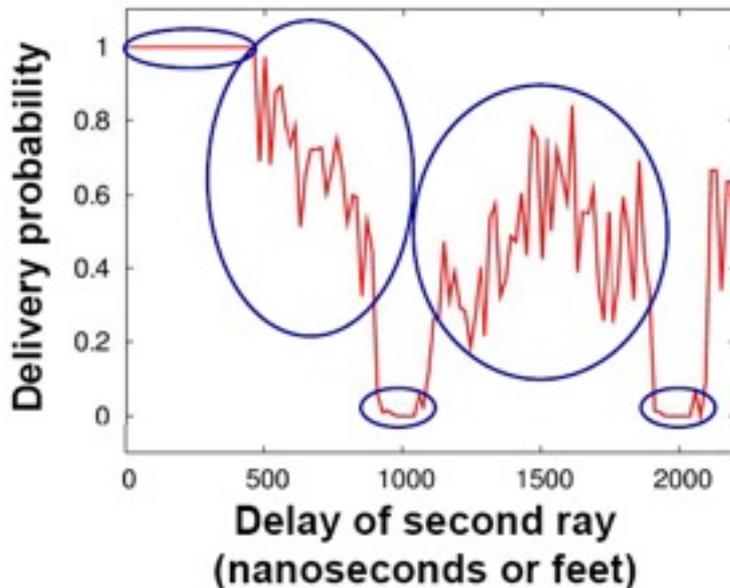
- Simulate with channel emulator



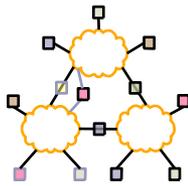
# A Plausible Explanation



- Multi-path can produce intermediate loss rates
- Appropriate multi-path delay is possible due to long-links

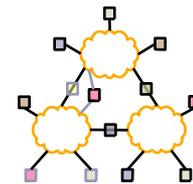


# Key Implications



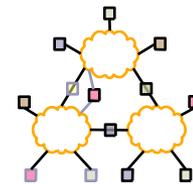
- Lack of a link abstraction!
  - Links aren't on or off... sometimes in-between
- Protocols must take advantage of these intermediate quality links to perform well
- How unique is this to Roofnet?
  - Cards designed for indoor environments used outdoors

# ETX measurement results



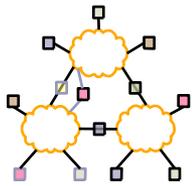
- Delivery *is* probabilistic
  - A  $1/r^2$  model wouldn't really predict this!
  - Sharp cutoff (by spec) of “good” vs “no” reception. Intermediate loss range band is just a few dB wide!
- Why?
  - Biggest factor: Multi-path interference
    - 802.11 receivers can suppress reflections  $< 250\text{ns}$
    - Outdoor reflections delay often  $> 1 \mu\text{sec}$
    - Delay offsets  $\Rightarrow$  symbol time look like valid symbols (large interference)
    - Offsets  $\neq$  symbol time look like random noise
    - Small changes in delay  $\Rightarrow$  big changes in loss rate

# Deciding Between Links



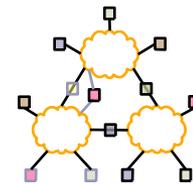
- Most early protocols: Hop Count
  - Link-layer retransmission can mask some loss
  - But: a 50% loss rate means your link is only 50% as fast!
- Threshold?
  - Can sacrifice connectivity. ☹️
  - Isn't a 90% path better than an 80% path?
- Real life goal: Find highest throughput paths

# Is there a better metric?



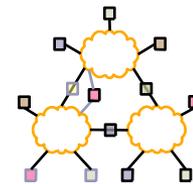
- Cut-off threshold
  - Disconnected network
- Product of link delivery ratio along path
  - Does not account for inter-hop interference
- Bottleneck link (highest-loss-ratio link)
  - Same as above
- End-to-end delay
  - Depends on interface queue lengths

# ETX Metric Design Goals



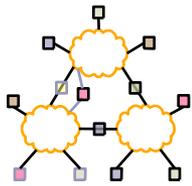
- Find high throughput paths
- Account for lossy links
- Account for asymmetric links
- Account for inter-link interference
- Independent of network load (don't incorporate congestion)

# Forwarding Packets is Expensive



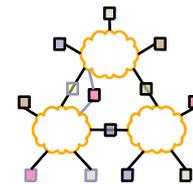
- Throughput of 802.11b  $\approx$  11Mbits/s
  - In reality, you can get about 5.
- What is throughput of a chain?
  - $A \rightarrow B \rightarrow C$  ?
  - $A \rightarrow B \rightarrow C \rightarrow D$  ?
  - Assume minimum power for radios.
- Routing metric should take this into account! Affects throughput

# ETX



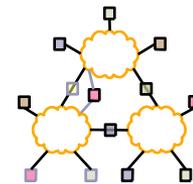
- Measure each link's delivery probability with broadcast probes (& measure reverse)
- $P(\text{delivery}) = (d_f * d_r)$  (ACK must be delivered too...)
- Link ETX =  $1 / P(\text{delivery})$
- Route ETX =  $\Sigma$  link ETX
  - Assumes all hops interfere - not true, but seems to work okay so far

# ETX: Sanity Checks



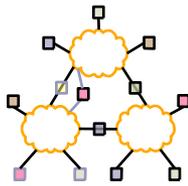
- ETX of perfect 1-hop path: 1
- ETX of 50% delivery 1-hop path: 2
- ETX of perfect 3-hop path: 3
  
- (So, e.g., a 50% loss path is better than a perfect 3-hop path! A threshold would probably fail here...)

# Rate Adaptation



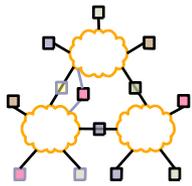
- What if links @ different rates?
- ETT – expected *transmission time*
  - $ETX / \text{Link rate} = 1 / ( P(\text{delivery}) * \text{Rate} )$
- What is best rate for link?
  - The one that minimizes ETT for the link!
  - SampleRate is a technique to adaptively figure this out.

# Discussion



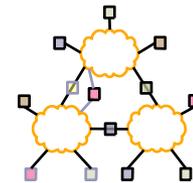
- Value of implementation & measurement
  - Simulators did not “do” multipath
    - Routing protocols dealt with the simulation environment just fine
    - Real world behaved differently and really broke a lot of the proposed protocols that worked so well in simulation!
- Rehash: Wireless differs from wired...
- Metrics: Optimize what matters; hop count often a very bad proxy in wireless
- What we didn't look at: routing protocol overhead
  - One cool area: Geographic routing

# Overview

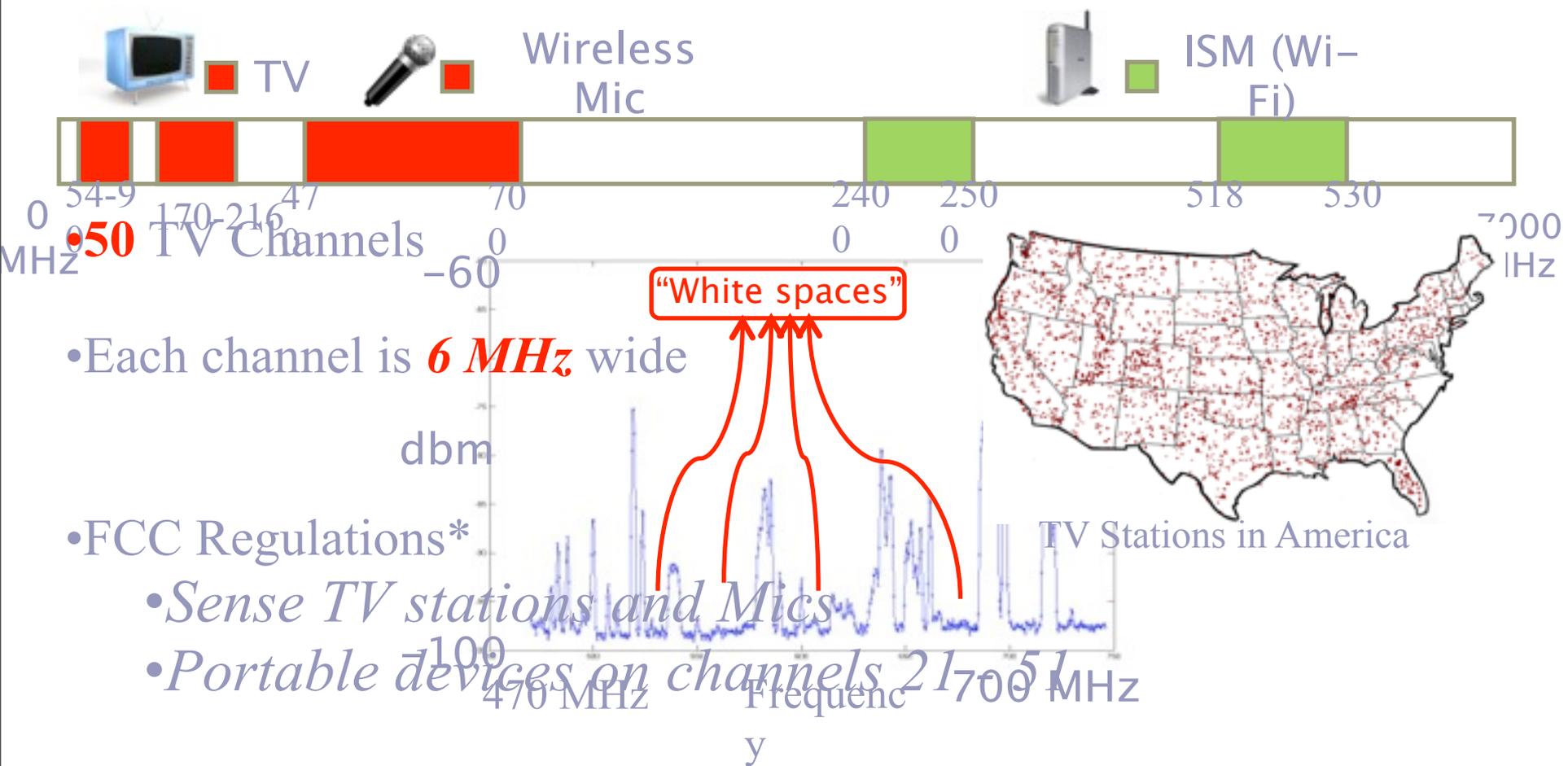


- 802.11
  - Deployment patterns
  - Reaction to interference
  - Interference mitigation
- Mesh networks
  - Architecture
  - Measurements
- **White space networks**





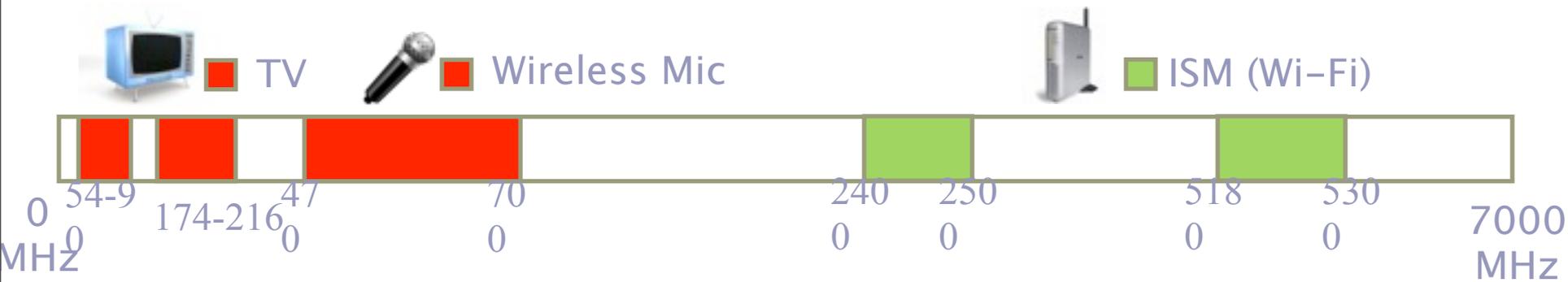
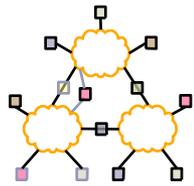
# What are White Spaces?



**White Spaces**

are *Unoccupied* TV Channels

# The Promise of White Spaces



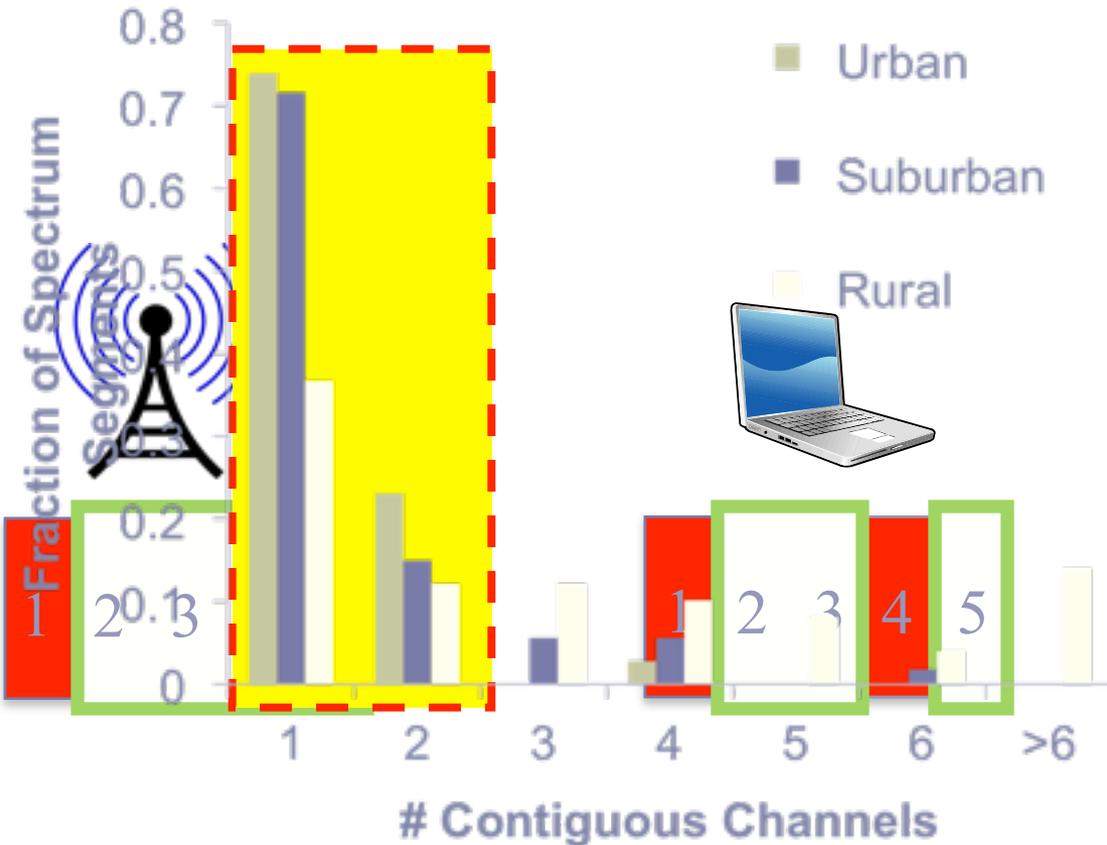
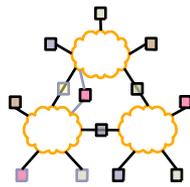
Up to 3x of 802.11g

**More  
Spectrum**

**Longer  
Range**

*at least 3 - 4x of Wi-Fi*

# White Spaces Spectrum Availability



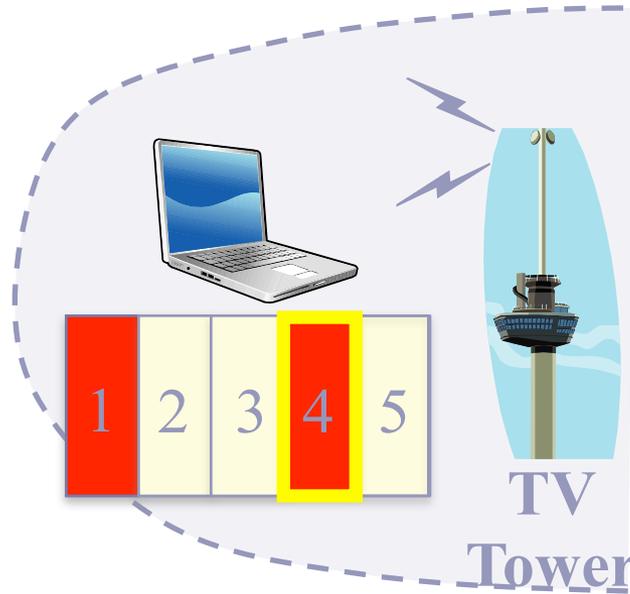
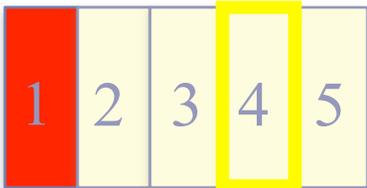
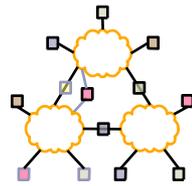
## Differences from ISM(Wi-Fi)

**Fragmentation**

Variable channel widths

Each TV Channel is 6 MHz    Spectrum is Fragmented    Channels for more bandwidth

# White Spaces Spectrum Availability



## Differences from ISM(Wi-Fi)

### Fragmentation

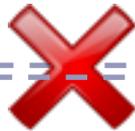
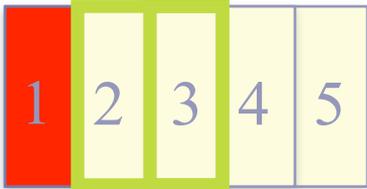
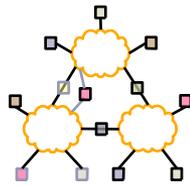
↳ Variable channel widths

### Spatial Variation

↳ Cannot assume same channel free everywhere

Location impacts spectrum availability ⇒ Spectrum exhibits spatial variation

# White Spaces Spectrum Availability



## Differences from ISM(Wi-Fi)

### Fragmentation

↳ Variable channel widths

### Spatial Variation

↳ Cannot assume same channel free everywhere

### Temporal Variation

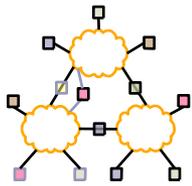
↳ Same Channel will not always be free

↳ Any connection can be disrupted *any* time

Incumbents appear/disappear over time

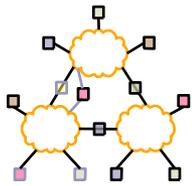
⇒ Must reconfigure after disconnection

# Channel Assignment in Wi-Fi



Fixed Width Channels  $\Rightarrow$  Optimize *which* channel to use

# Spectrum Assignment in WhiteFi



## Spectrum Assignment Problem

Goal

Maximize Throughput

Include

Spectrum at clients

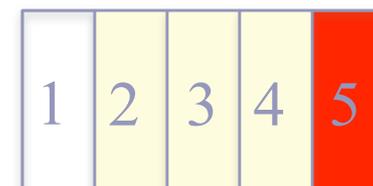
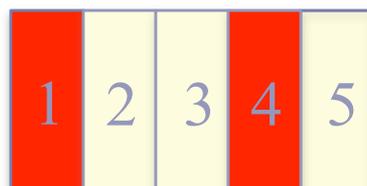
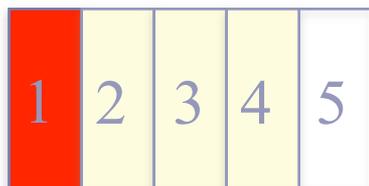
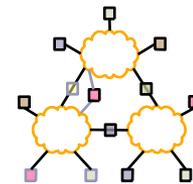
Assign

Center Channel  
&  
Width

Fragmentation  $\Rightarrow$  Optimize for *both*, center channel and width

Spatial Variation  $\Rightarrow$  BS must use channel *iff* free at client

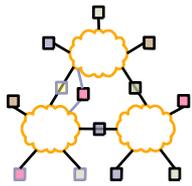
# Accounting for Spatial Variation



$$\begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline \end{array} \cup \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline \end{array} \cup \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline \end{array} = \begin{array}{|c|c|c|c|c|} \hline 1 & 2 & 3 & 4 & 5 \\ \hline \end{array}$$

The equation illustrates the union of three sets. Each set is a 5-element array. The first set has the first element (1) highlighted in red. The second set has the fourth element (4) highlighted in red. The third set has the fifth element (5) highlighted in red. The union of these three sets is shown as a single 5-element array where all elements (1, 2, 3, 4, 5) are highlighted in red, and the entire array is enclosed in a yellow box.

# Intuition

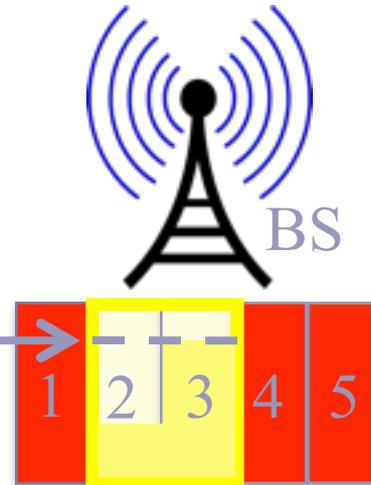


## Intuition

Use widest possible channel

But

Limited by *most* busy channel



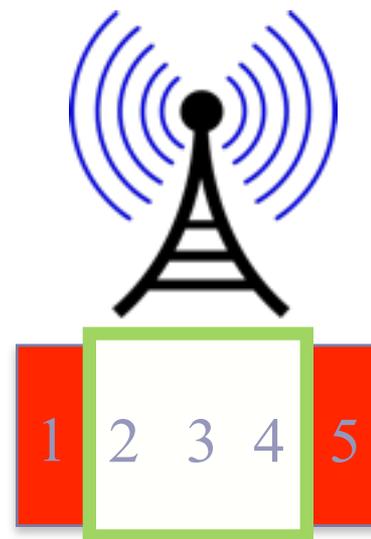
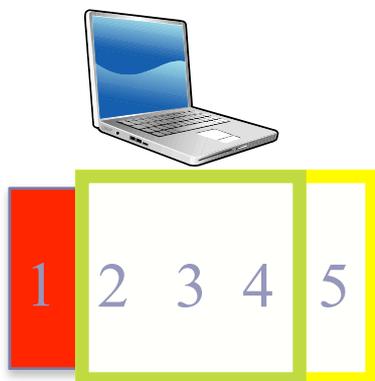
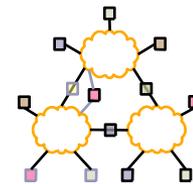
▪ Carrier Sense Across **All** Channels

▪ **All** channels must be free

$$\rho_{\text{BS}}(2 \text{ and } 3 \text{ are free}) = \rho_{\text{BS}}(2 \text{ is free}) \times \rho_{\text{BS}}(3 \text{ is free})$$

**Tradeoff between wider channel widths  
and opportunity to transmit on each channel**

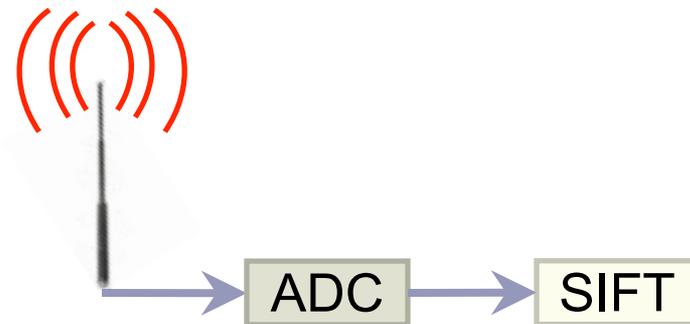
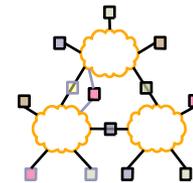
# Discovering a Base Station



**Discovery Time =  $O(B \times W)$**

Fragmentation  $\Rightarrow$  Try different center channel and widths  
channels used by the BS?

# SIFT, by example



## SIFT

Does *not* decode packets

Pattern match in time domain

